## Vogel Approximation Method (VAM)

VAM is an improved version of the least cost method that generally produces better solutions. The steps involved in this method are:

Step 1: For each row (column) with strictly positive capacity (requirement), determine a penalty by subtracting the smallest unit cost element in the row (column) from the next smallest unit cost element in the same row (column).
Step 2: Identify the row or column with the largest penalty among all the rows and columns. If the penalties corresponding to two or more rows or columns are equal we select the topmost row and the extreme left column.

Step 3: We select $\mathrm{X}_{\mathrm{ij}}$ as a basic variable if $\mathrm{C}_{\mathrm{ij}}$ is the minimum cost in the row or column with largest penalty. We choose the numerical value of $\mathrm{X}_{\mathrm{ij}}$ as high as possible subject to the row and the column constraints. Depending upon whether $a_{i}$ or $b_{j}$ is the smaller of the two $\mathrm{i}^{\text {th }}$ row or $\mathrm{j}^{\text {th }}$ column is crossed out.

Step 4: The Step 2 is now performed on the uncrossed-out rows and columns until all the basic variables have been satisfied.

## Example

Consider the following transportation problem

| Origin | Destination |  |  |  | ai |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 | 2 | 3 | 4 |  |
| 1 | 20 | 22 | 17 | 4 | 70 |
| 2 | 24 | 37 | 9 | 7 | 50 |
| 3 | 32 | 37 | 20 | 15 | 110 |
| $\mathrm{~b}_{\mathrm{j}}$ | 60 | 40 | 30 | 240 |  |

Note: $\mathrm{a}_{\mathrm{i}}=$ capacity (supply)
$\mathrm{b}_{\mathrm{j}}=$ requirement (demand)
Now, compute the penalty for various rows and columns which is shown in the following table:

| Origin | Destination |  |  |  |  | $\mathrm{a}_{\mathrm{i}}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 | 2 | 3 | 4 |  |  |
| Column |  |  |  |  |  |  |
| Penalty |  |  |  |  |  |  |$]$

Look for the highest penalty in the row or column, the highest penalty occurs in the second column and the minimum unit cost i.e. $\mathrm{c}_{\mathrm{ij}}$ in this column is $\mathrm{c}_{12}=22$. Hence assign 40 to this cell i.e. $x_{12}=40$ and cross out the second column (since second column was satisfied). This is shown in the following table:

| Origin | Destination |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 | 2 | 3 | 4 | Column <br> Penalty |  |
| 1 | 20 | 22 | $\mathbf{4 0}$ | 17 | 4 | 80 |
| 2 | 24 | 37 | 9 | 7 | 70 | 2 |
| 3 | 32 | 37 | 20 | 15 | 50 | 5 |
| $\mathrm{~b}_{\mathrm{j}}$ | 60 | 40 | 30 | 110 | 240 |  |
| Row Penalty | 4 | 15 |  | 8 | 3 |  |

The next highest penalty in the uncrossed-out rows and columns is 13 which occur in the first row and the minimum unit cost in this row is $\mathrm{c}_{14}=4$, hence $\mathrm{x}_{14}=80$ and cross out the first row. The modified table is as follows:

\left.| Origin | Destination |  |  |  |  | ai |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 | 2 | 3 | 4 |  |  |
| Column |  |  |  |  |  |  |
| Penalty |  |  |  |  |  |  |$\right]$

The next highest penalty in the uncrossed-out rows and columns is 8 which occurs in the third column and the minimum cost in this column is $\mathrm{c}_{23}=9$, hence $\mathrm{x}_{23}=30$ and cross out the third column with adjusted capacity, requirement and penalty values.
The modified table is as follows:

| Origin | Destination |  |  |  |  | $\mathrm{a}_{\mathrm{i}}$ | Column <br> Penalty |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 | 2 |  | 3 | 4 |  |  |
| 1 | 20 | 22 |  | 17 |  | 0 | 13 |
|  |  |  | 40 |  |  |  |  |
| 2 | 24 | 37 |  | 9 | 7 | 40 | 17 |
| 3 | 32 | 37 |  | 20 | 15 | 50 | 17 |
| $\mathrm{b}_{\mathrm{j}}$ | 60 | $4{ }^{4}$ |  | 30 | 110 | 240 |  |
| Row Penalty | 8 | 15 |  | 8 | 8 |  |  |

The next highest penalty in the uncrossed-out rows and columns is 17 which occurs in the second row and the smallest cost in this row is $\mathrm{c}_{24}=15$, hence $\mathrm{x}_{24}=30$ and cross out the fourth column with the adjusted capacity, requirement and penalty values. The modified table is as follows:


The next highest penalty in the uncrossed-out rows and columns is 17 which occurs in the second row and the smallest cost in this row is $c_{21}=24$, hence $\mathrm{xi}_{21}=10$ and cross out the second row with the adjusted capacity, requirement and penalty values.
The modified table is as follows:

| Origin | Destination |  |  |  |  |  | $\mathrm{a}_{\mathrm{i}}$ | Column Penalty |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 | 2 |  | 3 | 4 |  |  |  |
| 1 | 20 |  | 40 | 17 |  | 4 80 | 0 | 13 |
| 2 | 24 | 37 |  | 9 | 7 |  | 0 | 17 |
|  | 10 |  |  |  |  | 30 |  |  |
| 3 | 32 | 37 |  | 20 | 15 |  | 50 | 17 |
| $\mathrm{b}_{\mathrm{j}}$ | 60 | $4 \downarrow$ |  | 30 | 110 |  | 240 |  |
| Row Penalty | 8 | 15 |  | 8 | 8 |  |  |  |

The next highest penalty in the uncrossed-out rows and columns is 17 which occurs in the third row and the smallest cost in this row is $c_{31}=32$, hence $\mathrm{xi}_{31}=50$ and cross out the third row or first column.
The modified table is as follows:

| Origin | Destination |  |  |  |  |  | $\mathrm{a}_{\mathrm{i}}$ | Column Penalty |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 1 | 2 |  | 3 | 4 |  |  |  |
| 1 | 20 | 22 |  | 17 |  | 4 | 0 | 13 |
|  |  |  | 40 |  | , | 80 |  |  |
| 2 | 24 | 37 |  | 9 |  |  | 0 | 17 |
|  | 10 |  |  |  |  | 30 |  |  |
| 3 | 32 | 37 |  | 20 | 15 |  | 0 | 17 |
|  | 50 |  |  |  |  |  |  |  |
|  | 60 | 40 |  | 30 | 110 |  | 24 |  |
| $\mathrm{b}_{\mathrm{j}}$ |  |  |  |  |  |  |  |  |
| Row Penalty | 8 | 15 |  | 8 | 8 |  |  |  |

The transportation cost corresponding to this choice of basic variables is
$22 * 40+4 * 80+9 * 30+7 * 30+24 * 10+32 * 50=3520$

